

# On $r$ - dynamic coloring of the gear graph families 

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#### Abstract

:

An r-dynamic coloring of a graph $G$ is a proper coloring $c$ of the vertices such that $|c(N(v))| \geq \min \{r, d(v)\}$, for each $v \in V(G)$. The $r$-dynamic chromatic number of a graph $G$ is the minimum $k$ such that $G$ has an r-dynamic coloring with $k$ colors. In this paper, we obtain the $r$-dynamic chromatic number of the middle, central and line graphs of the gear graph.


Keywords: $r$ - dynamic coloring; gear graph; middle graph; central graph and line graph.

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## 1. Introduction

Graphs in this paper are simple and finite. For undefined terminologies and notations see [5, 17]. Thus for a graph $G, \delta(G), \Delta(G)$ and $\chi(G)$ denote the minimum degree, maximum degree and chromatic number of $G$ respectively. When the context is clear we write, $\delta, \Delta$ and $\chi$ for brevity. For $v \in V(G)$, let $N(v)$ denote the set of vertices adjacent to $v$ in $G$ and $d(v)=|N(v)|$. The $r$-dynamic chromatic number was first introduced by Montgomery [14].

An $r$-dynamic coloring of a graph $G$ is a map $c$ from $V(G)$ to the set of colors such that (i) if $u v \in E(G)$, then $c(u) \neq c(v)$ and (ii) for each vertex $v \in V(G),|c(N(v))| \geq \min \{r, d(v)\}$, where $N(v)$ denotes the set of vertices adjacent to $v$ and $d(v)$ its degree and $r$ is a positive integer.

The first condition characterizes proper colorings, the adjacency condition and second condition is double-adjacency condition. The $r$-dynamic chromatic number of a graph $G$, written $\chi_{r}(G)$, is the minimum $k$ such that $G$ has an $r$-dynamic proper $k$-coloring. The 1-dynamic chromatic number of a graph G is equal to its chromatic number. The 2-dynamic chromatic number of a graph has been studied under the name dynamic chromatic number denoted by $\chi_{d}(G)[1,2,3,4,8]$. By simple observation, we can show that $\chi_{r}(G) \leq \chi_{r+1}(G)$, however $\chi_{r+1}(G)-\chi_{r}(G)$ can be arbitrarily large, for example $\chi($ Petersen $)=2, \chi_{d}($ Petersen $)=3$, but $\chi_{3}($ Petersen $)=10$. Thus, finding an exact values of $\chi_{r}(G)$ is not trivially easy.

There are many upper bounds and lower bounds for $\chi_{d}(G)$ in terms of graph parameters. For example, for a graph $G$ with $\Delta(G) \geq 3$, Lai et al. [8] proved that $\chi_{d}(G) \leq \Delta(G)+1$. An upper bound for the dynamic chromatic number of a $d$-regular graph $G$ in terms of $\chi(G)$ and the independence number of $G, \alpha(G)$, was introduced in [7]. In fact, it was proved that $\chi_{d}(G) \leq \chi(G)+2 \log _{2} \alpha(G)+3$. Taherkhani gave in [15] an upper bound for $\chi_{2}(G)$ in terms of the chromatic number, the maximum degree $\Delta$ and the minimum degree $\delta$. i.e., $\chi_{2}(G)-\chi(G) \leq\left\lceil(\Delta e) / \delta \log \left(2 e\left(\Delta^{2}+1\right)\right)\right\rceil$.

Li et al.proved in [10] that the computational complexity of $\chi_{d}(G)$ for a 3-regular graph is an NP-complete problem. Furthermore, Li and Zhou [9] showed that to determine whether there exists a 3 -dynamic coloring, for a claw free graph with the maximum degree 3, is NP-complete.
N.Mohanapriya et al. [11, 12] studied the dynamic chromatic number for various graph families. Also, it was proven in [13] that the $r$-dynamic chromatic number of line graph of a helm graph $H_{n}$.

In this paper, we study $\chi_{r}(G)$, when $1 \leq r \leq \Delta$. We find the $r$ - dynamic chromatic number of the middle, central and line graphs of the gear graph.

## 2. Preliminaries

Let $G$ be a graph with vertex set $V(G)$ and edge set $E(G)$. The middle graph [6] of $G$, denoted by $M(G)$ is defined as follows. The vertex set of $M(G)$ is $V(G) \cup E(G)$. Two vertices $x, y$ of $M(G)$ are adjacent in $M(G)$ in case one of the following holds: (i) $x, y$ are in $E(G)$ and $x, y$ are adjacent in $G$. (ii) $x$ is in $V(G), y$ is in $E(G)$, and $x, y$ are incident in $G$.

The central graph [16] $C(G)$ of a graph $G$ is obtained from $G$ by adding an extra vertex on each edge of $G$, and then joining each pair of vertices of the original graph which were previously non-adjacent.

The line graph [13] of $G$ denoted by $L(G)$ is the graph with vertices are the edges of $G$ with two vertices of $L(G)$ adjacent whenever the corresponding edges of $G$ are adjacent.

The gear graph is a wheel graph with a graph vertex added between each pair of adjacent graph vertices of the outer cycle. The gear graph $G_{n}$ has $2 n+1$ nodes and $3 n$ edges.

Let $V\left(G_{n}\right)=\{v\} \cup\left\{v_{i}: 1 \leq i \leq n\right\} \cup\left\{u_{i}: 1 \leq i \leq n\right\}$ and $E\left(G_{n}\right)=\left\{v v_{i}: 1 \leq i \leq n\right\} \cup\left\{u_{i} v_{i}: 1 \leq i \leq n\right\} \cup\left\{u_{i} v_{i+1}: 1 \leq i \leq n\right.$ and the meaning of $\bmod \mathrm{n}$ is the obvious $\}$.

## 3. Main Theorem

Theorem 3.1. Let $n \geq 5, M\left(G_{n}\right)$ be the middle graph of a gear graph $G_{n}$ and let
$\Delta=\Delta\left(M\left(G_{n}\right)\right)$. Then

$$
\begin{aligned}
& n+1, \quad 1 \leq r \leq 4 \\
& n+2, \quad 5 \leq r \leq \Delta-2 \\
& n+4, \quad r=\Delta-1 \text { and } n \equiv 0 \bmod 3 \\
& \chi_{r}\left(M\left(G_{n}\right)\right)=\begin{array}{lll}
n+5, & r=\Delta-1 & \text { and } n \equiv 1 \bmod 3 \\
n+4, & r=\Delta-1 & \text { and } n \equiv 2
\end{array} \bmod 3 \\
& n+5, \quad r=\Delta \text { and } n \equiv 0 \bmod 3 \\
& n+7, \quad r=\Delta \text { and } n \equiv 1 \bmod 3 \\
& n+6, \quad r=\Delta \text { and } n \equiv 2 \bmod 3
\end{aligned}
$$

Proof. By the definition of middle graph,
$V\left(M\left(G_{n}\right)\right)=V\left(G_{n}\right) \cup E\left(G_{n}\right)=\{v\} \cup\left\{v_{i}: 1 \leq i \leq n\right\} \cup\left\{u_{i}: 1 \leq i \leq n\right\} \cup$ $\left\{e_{i}: 1 \leq i \leq n\right\} \cup\left\{s_{i}: 1 \leq i \leq 2 n\right\}$.

The vertices $v$ and $\left\{e_{i}: 1 \leq i \leq n\right\}$ induces a clique of order $K_{n+1}$ in $M\left(G_{n}\right)$.

Thus, $\chi_{\delta}\left(M\left(G_{n}\right)\right) \geq n+1$.
We divide the proof into some cases.

## Case 1 : For $1 \leq r \leq 4$

The $r$ - dynamic $(n+1)$ coloring is as follows:
For $1 \leq i \leq n$, assign the color $c_{i}$ to $e_{i}$ and assign the color $c_{n+1}$ to $v$.
For $1 \leq i \leq n$, assign the color $c_{n+1}$ to $u_{i}$ and $v_{i}$.
$\left|N\left(u_{i}\right)\right|=d\left(u_{i}\right)=2=\delta$,
$\left|N\left(v_{i}\right)\right|=d\left(v_{i}\right)=3$
$|N(v)|=d(v)=n$,
$\left|N\left(e_{i}\right)\right|=d\left(e_{i}\right)=n+3$
and $\left|N\left(s_{i}\right)\right|=d\left(s_{i}\right)=5$
For $1 \leq i \leq 2 n$, assign the allowed colors to the vertex $s_{i}$ and also it must satisfies the $r$-adjacency condition.

- color the vertices $s_{1}, s_{3}, s_{5}, s_{7}, \cdots s_{2 n-5}, s_{2 n-3}, s_{2 n-1}$ with colors $c_{3}, c_{4}, c_{5}, \cdots c_{n}, c_{1}, c_{2}$ (the order of assigned color is important).
- color the vertices $s_{2}, s_{4}, s_{6}, s_{8}, \cdots s_{2 n-4}, s_{2 n-2}, s_{2 n}$ with colors $c_{n}, c_{1}, c_{2}, c_{3}, \cdots c_{n-3}, c_{n-2}, c_{n-1}$ (the order of assigned color is important).

We know that the $|N(v)|=d(v)=n$, so we need the color $n+1$.
It is easy to verify that adjacency and r-adjacency conditions are fulfilled.

Hence, $\chi_{r}\left(M\left(G_{n}\right)\right)=n+1$, for $n \geq 5$ and $1 \leq r \leq 4$.

Case 2: For $5 \leq r \leq \Delta-2$
The $r$-dynamic $(n+2)$ coloring is as follows:
For $1 \leq i \leq n$, assign the color $c_{i}$ to $e_{i}$ and assign the color $c_{n+1}$ to $v$.
For $1 \leq i \leq n$, assign the color $c_{n+1}$ to $u_{i}$.
For $1 \leq i \leq 2 n$, if any, assign the vertex $s_{i}$ to one of the allowed colors

- such color exists, because $\left|N\left(s_{i}\right)\right|=d\left(s_{i}\right)=5$
- color the vertices $s_{1}, s_{3}, s_{5}, s_{7}, \cdots s_{2 n-5}, s_{2 n-3}, s_{2 n-1}$ with colors $c_{3}, c_{4}, c_{5}, \cdots c_{n}, c_{1}, c_{2}$ (the order of assigned color is important).
- color the vertices $s_{2}, s_{4}, s_{6}, s_{8}, \cdots s_{2 n-4}, s_{2 n-2}, s_{2 n}$ with colors $c_{n}, c_{1}, c_{2}, c_{3}, \cdots c_{n-3}, c_{n-2}, c_{n-1}$ (the order of assigned color is important).
- color the vertex $v_{i}$ with the color $c_{n+2}$.

Now $\left|N\left(s_{i}\right)\right|$ satisfies the $r$-adjacency condition.
But $d\left(e_{i}\right)=n+3$, so $N\left(e_{i}\right)$ having $n+2$ colors.
It is easy to verify that the $r$-adjacency condition is fulfilled.
Hence, $\chi_{r}\left(M\left(G_{n}\right)\right)=n+2$, for $n \geq 5$ and $5 \leq r \leq \Delta-2$.

Case 3: For $r=\Delta-1$ and $n \equiv 0 \bmod 3$
The $r$ - dynamic $(n+4)$ coloring is as follows:
For $1 \leq i \leq n$, assign the color $c_{i}$ to $e_{i}$ and assign the color $c_{n+1}$ to $v$.
For $1 \leq i \leq n$, assign the color $c_{n+1}$ to $u_{i}$.
For $1 \leq i \leq n$, assign the color $c_{n+2}$ to $v_{i}$.
$\left|N\left(e_{i}\right)\right|$ having $n+1$ colors only. So we assign one new color to $s_{i}$.

- color the vertices $s_{1}, s_{4}, s_{7}, s_{10}, \cdots s_{2 n-5}, s_{2 n-2}$ with color $c_{n+3}$.
- color the vertices $s_{2}, s_{5}, s_{8}, s_{11}, \cdots s_{2 n-4}, s_{2 n-1}$ with colors $c_{n+4}$.

Now $s_{3}, s_{6}, s_{9}, \cdots s_{2 n-3}, s_{2 n}$ are uncolored. So assign these vertices to any one of the allowed colors-such color exists.

- color the vertices $s_{3}, s_{6}, s_{9}, \cdots s_{2 n-3}, s_{2 n}$ with colors $c_{5}, c_{7}, c_{9}, \cdots c_{n-1}, c_{1}, c_{3}$ (the order of assigned color is important).

Now neighbours of $e_{i}$ having $n+4$ colors and an easy check shows that the $r$-adjacency condition is fulfilled.

Hence, $\chi_{r}\left(M\left(G_{n}\right)\right)=n+4$, for $n \geq 5, r=\Delta-1$ and $n \equiv 0 \bmod 3$.

Case 4 : For $r=\Delta-1$ and $n \equiv 1 \bmod 3$
The $r$ - dynamic $(n+5)$ coloring is as follows:
For $1 \leq i \leq n$, assign the color $c_{i}$ to $e_{i}$ and assign the color $c_{n+1}$ to $v$.
For $1 \leq i \leq n$, assign the color $c_{n+1}$ to $u_{i}$.
For $1 \leq i \leq n$, assign the color $c_{n+2}$ to $v_{i}$.
$N\left(e_{i}\right)$ having $n+1$ colors. So we have to assign one new color to $s_{i}$.

- color the vertices $s_{1}, s_{4}, s_{7}, s_{10}, \cdots s_{2 n-4}$ with color $c_{n+3}$.
- color the vertices $s_{2}, s_{5}, s_{8}, s_{1} 1, \cdots s_{2 n-3}$ with color $c_{n+4}$

But neighbours of $e_{n}$ having $n+1$ colors only. So we have to assign a new color $c_{n+5}$ to $s_{2 n-2}$.

Now neighbours of $e_{i}$ having $n+2$ colors. But the vertices $s_{2 n-1}$ and $s_{2 n}$ are uncolored.

So we have to assign any one of the allowed colors to $s_{2 n-1}$ and $s_{2 n}$.

- color the vertex $s_{2 n-1}$ with the color $c_{2}$ and color the vertex $s_{2 n}$ with the color $c_{3}$.

Now an easy check shows that the r-adjacency condition is fulfilled.
Hence, $\chi_{r}\left(M\left(G_{n}\right)\right)=n+5$, for $n \geq 5$ and $r=\Delta-1$ and $n \equiv 1 \bmod 3$.

Case 5 : For $r=\Delta-1$ and $n \equiv 2 \bmod 3$
The $r$ - dynamic $(n+4)$ coloring is as follows:
For $1 \leq i \leq n$, assign the color $c_{i}$ to $e_{i}$ and assign the color $c_{n+1}$ to $v$.
For $1 \leq i \leq n$, assign the color $c_{n+1}$ to $u_{i}$.
For $1 \leq i \leq n$, assign the color $c_{n+2}$ to $v_{i}$.

Now $N\left(e_{i}\right)$ having $n+1$ colors.so we have to assign one new color to $s_{i}$.

- color the vertices $s_{1}, s_{4}, s_{7}, s_{10}, \cdots s_{2 n-3}$ with color $c_{n+3}$.
- color the vertices $s_{2}, s_{5}, s_{8}, s_{1} 1, \cdots s_{2 n-2}$ with color $c_{n+4}$

But the vertices $s_{3}, s_{6}, s_{9}, \cdots s_{2 n-4}, s_{2 n-1}$ and $s_{2 n}$ are uncolored. So we have to assign any one of the allowed colors to these vertices.

- color the vertices $s_{3}, s_{6}, s_{9}, \cdots s_{2 n-1}, s_{2 n}$ with colors $c_{4}, c_{1} 1, c_{7}, c_{3}, \cdots, c_{2}, c_{8}$ respectively.(the order of assigned color is important).

Now an easy check shows that the $r$ - adjacency condition is fulfilled. Hence, $\chi_{r}\left(M\left(G_{n}\right)\right)=n+4$, for $n \geq 5, r=\Delta-1$ and $n \equiv 2 \bmod 3$.

Case 6 : For $r=\Delta$ and $n \equiv 0 \bmod 3$
The $r$ - dynamic $(n+5)$ coloring is as follows:
For $1 \leq i \leq n$, assign the color $c_{i}$ to $e_{i}$ and assign the color $c_{n+1}$ to $v$.
For $1 \leq i \leq n$, assign the color $c_{n+1}$ to $u_{i}$.
For $1 \leq i \leq n$, assign the color $c_{n+2}$ to $v_{i}$.
For $r=\Delta$, we have to assign two new colors to neighbours of $e_{i}$.

- color the vertices $s_{1}, s_{4}, s_{7}, s_{10}, \cdots s_{2 n-5}, s_{2 n-2}$ with color $c_{n+3}$
- color the vertices $s_{2 n}, s_{3}, s_{6}, s_{9}, \cdots s_{2 n-3}$ with color $c_{n+4}$
- color the vertices $s_{2}, s_{5}, s_{8}, s_{11}, \cdots s_{2 n-1}$ with color $c_{n+5}$

Now an easy check shows that the $r$-adjacency condition is fulfilled for all the vertices.

Hence, $\chi_{r}\left(M\left(G_{n}\right)\right)=n+5$, for $n \geq 5, r=\Delta$ and $n \equiv 0 \bmod 3$.

Case 7 : For $r=\Delta$ and $n \equiv 1 \bmod 3$
The $r$ - dynamic $(n+7)$ coloring is as follows:
For $1 \leq i \leq n$, assign the color $c_{i}$ to $e_{i}$ and assign the color $c_{n+1}$ to $v$.
For $1 \leq i \leq n$, assign the color $c_{n+1}$ to $u_{i}$.
For $1 \leq i \leq n$, assign the color $c_{n+2}$ to $v_{i}$.
For $r=\Delta$, we have to assign two new colors to neighbours of $e_{i}$.

- color the vertices $s_{1}, s_{4}, s_{7}, s_{10}, \cdots s_{2 n-4}$ with color $c_{n+3}$.
- color the vertices $s_{2 n}, s_{3}, s_{6}, s_{9}, \cdots s_{2 n-5}$ with color $c_{n+4}$.
- color the vertices $s_{2}, s_{5}, s_{8}, s_{11}, \cdots s_{2 n-3}$ with color $c_{n+5}$.

But neighbours of $e_{n}$ does not satisfies the $r$-adjacency condition.
So we have to assign two new colors to the vertices $s_{2 n-2}$ and $s_{2 n-1}$ respectively.

- color the vertex $s_{2 n-2}$ with the color $c_{n+6}$ and color the vertex $s_{2 n-1}$ with the color $c_{n+7}$.

So we have to assign any one of the allowed colors to $s_{2 n-1}$ and $s_{2 n}$.
Now an easy check shows that the r-adjacency condition is fulfilled.
Hence, $\chi_{r}\left(M\left(G_{n}\right)\right)=n+7$, for $n \geq 5, r=\Delta$ and $n \equiv 1 \bmod 3$.

Case 8 : For $r=\Delta$ and $n \equiv 2 \bmod 3$
The $r$ - dynamic $(n+6)$ coloring is as follows:
For $1 \leq i \leq n$, assign the color $c_{i}$ to $e_{i}$ and assign the color $c_{n+1}$ to $v$.
For $1 \leq i \leq n$, assign the color $c_{n+1}$ to $u_{i}$.
For $1 \leq i \leq n$, assign the color $c_{n+2}$ to $v_{i}$.
For $r=\Delta$, we have to assign two new colors to neighbours of $e_{i}$.

- color the vertices $s_{1}, s_{4}, s_{7}, s_{10}, \cdots s_{2 n-3}$ with color $c_{n+3}$.
- color the vertices $s_{2 n}, s_{3}, s_{6}, s_{9}, \cdots s_{2 n-4}$ with color $c_{n+4}$.
- color the vertices $s_{2}, s_{5}, s_{8}, s_{11}, \cdots s_{2 n-2}$ with color $c_{n+5}$.

Now neighbours of $e_{n}$ does not satisfies the $r$ - adjacency condition.

- color the vertex $s_{2 n-1}$ with the new color $c_{n+6}$

Now an easy check shows that the $r-$ adjacency condition is fulfilled. Hence, $\chi_{r}\left(M\left(G_{n}\right)\right)=n+6$, for $n \geq 5, r=\Delta$ and $n \equiv 2 \bmod 3$.

Theorem 3.2. Let $n \geq 5, C\left(G_{n}\right)$ be the central graph of a Gear graph $G_{n}$ and let
$\Delta=\Delta\left(C\left(G_{n}\right)\right)$. Then

$$
\chi_{r}\left(C\left(G_{n}\right)\right)=\left\{\begin{array}{l}
n+1, \quad r=1 \\
2 n+1, \quad \delta \leq r \leq \Delta-2 \\
2 n+2, \quad r=\Delta-1 \\
3 n+3, \quad r=\Delta
\end{array}\right.
$$

Proof. By the definition of central graph, subdividing each edge of $G_{n}$ exactly once and then joining each pair of vertices of $G_{n}$ which were non-adjacent.

Let $V\left(C\left(G_{n}\right)\right)=V\left(G_{n}\right) \cup E\left(G_{n}\right)=\{v\} \cup\left\{v_{i}: 1 \leq i \leq n\right\} \cup\left\{u_{i}: 1 \leq i \leq n\right\} \cup$ $\left\{e_{i}: 1 \leq i \leq n\right\} \cup\left\{s_{i}: 1 \leq i \leq 2 n\right\}$

We divide the proof into some cases.
Case 1: For $r=1$
The $r$ - dynamic $(n+1)-$ coloring is as follows:
For $1 \leq i \leq n$, assign the color $c_{i}$ to $v_{i}$ and $u_{i}$.
For $1 \leq i \leq n-1$, assign the color $c_{i}$ to $e_{i+1}$ and assign the color $c_{n}$ to $e_{1}$.
$\left|N\left(u_{i}\right)\right|=d\left(u_{i}\right)=2 n$
$\left|N\left(v_{i}\right)\right|=d\left(v_{i}\right)=2 n$
$|N(v)|=d(v)=2 n$,
$\left|N\left(e_{i}\right)\right|=d\left(e_{i}\right)=2$
and $\left|N\left(s_{i}\right)\right|=d\left(s_{i}\right)=2$
For $1 \leq i \leq 2 n$, assign the color $c_{n+1}$ to the vertex $s_{i}$ and assign the color $c_{n+1}$ to $v$.

Now an easy check shows that the $r$ - adjacency condition is fulfilled.
Hence, $\chi_{r}\left(C\left(G_{n}\right)\right)=n+1$, for $r=1$.

Case 2: For $\delta \leq r \leq \Delta-2$
The $r$ - dynamic $(2 n+1)-$ coloring is as follows:
For $1 \leq i \leq n$, assign the color $c_{i}$ to $v_{i}$.
For $1 \leq i \leq 2 n$, assign the color $c_{n+1}$ to $s_{i}$.
For $1 \leq i \leq n-1$, assign the color $c_{i}$ to $e_{i+1}$ and assign the color $c_{n}$ to $e_{1}$ and also assign the color $c_{n+1}$ to $v$.

- Color the vertices $u_{1}, u_{2}, u_{3}, \cdots u_{n-1}, u_{n}$ with colors $c_{n+2}, c_{n+3}, \cdots c_{2 n}, c_{2 n+1}$ (the order of assigned color is important).

Now an easy check shows that the $r$-adjacency condition is fulfilled.
Hence, $\chi_{r}\left(C\left(G_{n}\right)\right)=2 n+1$, for $\delta \leq r \leq \Delta-2$

Case 3 : For $r=\Delta-1$
The $r$ - dynamic $(2 n+2)-$ coloring is as follows:
For $1 \leq i \leq n$, assign the color $c_{i}$ to $v_{i}$ and assign the color $c_{n+1}$ to $v$.
For $1 \leq i \leq n-1$, assign the color $c_{i}$ to $e_{i+1}$ and assign the color $c_{n}$ to $e_{1}$.

- Color the vertices $u_{1}, u_{2}, u_{3}, \cdots u_{n-1}, u_{n}$ with colors $c_{n+2}, c_{n+3}, \cdots c_{2 n}, c_{2 n+1}$ (the order of assigned color is important).
- Color the vertices $s_{2}, s_{4}, s_{6}, \cdots s_{2 n-2}, s_{2 n}$ with color $c_{n+1}$.
- Color the vertices $s_{1}, s_{3}, s_{5}, \cdots s_{2 n-3}, s_{2 n-1}$ with colors $c_{2 n+2}$.

Now an easy check shows that the $r$-adjacency condition is fulfilled.
Hence, $\chi_{r}\left(C\left(G_{n}\right)\right)=2 n+2$, for $r=\Delta-1$

Case 4 : For $r=\Delta$
The $r$ - dynamic $(3 n+3)-$ coloring is as follows:
For $1 \leq i \leq n$, assign the color $c_{i}$ to $v_{i}$ and assign the color $c_{n+1}$ to $v$.

- Color the vertices $u_{1}, u_{2}, u_{3}, \cdots u_{n-1}, u_{n}$ with colors $c_{n+2}, c_{n+3}, \cdots c_{2 n}, c_{2 n+1}$ (the order of assigned color is important).
- Color the vertices $s_{1}, s_{3}, s_{5}, \cdots s_{2 n-3}, s_{2 n-1}$ with colors $c_{2 n+2}$ and color the vertices $s_{2}, s_{4}, s_{6}, \cdots s_{2 n-2}, s_{2 n}$ with color $c_{2 n+3}$.
- Color the vertices $e_{1}, e_{2}, e_{3}, \cdots e_{n-1}, e_{n}$ with colors $c_{2 n+4}, c_{2 n+5}, c_{2 n+6} \cdots c_{3 n+2}, c_{3 n+3}$ respectively.(the order of assigned color is important).

Now an easy check shows that the $r$-adjacency condition is fulfilled.
Hence, $\chi_{r}\left(C\left(G_{n}\right)\right)=3 n+3$, for $r=\Delta$

## Result:

Let us consider the line graphs built on the base of Gear graph.
By the definition of line graph
$V\left(L\left(G_{n}\right)\right)=E\left(G_{n}\right)=\left\{e_{i}: 1 \leq i \leq n\right\} \cup\left\{s_{i}: 1 \leq i \leq 2 n\right\}$.
Note that $d\left(e_{i}\right)=n+1, d\left(s_{i}\right)=3$. Hence $\delta\left(L\left(G_{n}\right)\right)=3$.
Next, observe that the vertices $\left\{e_{1}, e_{2}, e_{3}, . ., e_{n}\right\}$ induces a clique $K_{n}$ in $L\left(G_{n}\right)$. Thus,

$$
\begin{equation*}
\chi_{\delta}\left(L\left(G_{n}\right)\right) \geq n \tag{3.1}
\end{equation*}
$$

for any $r$. Let us start with $r=\delta$.

Proposition 3.3. Let $n \geq 5$.Let $L\left(G_{n}\right)$ be the line graph of a Gear graph $G_{n}$. Then $\chi_{\delta}\left(L\left(G_{n}\right)\right)=n$.

Proof. Due to (1), we have $\chi_{\delta}\left(L\left(G_{n}\right)\right) \geq n$.
So, we need to fix only appropriate coloring.
For $1 \leq i \leq n$, assign the color $i$ to $e_{i}$. Next,assign the colors to $s_{i}$ such that partial coloring is proper and the $r$ - adjacency condition for $r=\delta$ is also fulfilled.

That is we should assign one of the allowed colors from $\{1,2, \cdots n\}$ to vertex $s_{i}$ of degree $3,1 \leq i \leq n$.

The coloring we obtained is $\delta-$ dynamic coloring of $L\left(G_{n}\right)$.
The result from proposition can be extended to $r$ - dynamic coloring for line graph of Gear graph for all $r$, where $1 \leq r \leq \Delta$.

Theorem 3.4. Let $n \geq 6, L\left(G_{n}\right)$ be the line graph of a Gear graph $G_{n}$ and
let $\Delta=\Delta\left(L\left(G_{n}\right)\right)$. Then

$$
\chi_{r}\left(L\left(G_{n}\right)\right)=\left\{\begin{array}{l}
n, \quad 1 \leq r \leq n-1 \\
n+2, \quad r=n \text { and } n \not \equiv 1 \bmod 3 \\
n+3, \quad r=n \quad \text { and } n \equiv 1 \bmod 3 \\
n+3, \quad r=n+1=\Delta, \quad n \geq 5 \text { and } 2 n \equiv 0 \bmod 3 \\
n+4, \quad r=n+1=\Delta, \quad n \geq 5 \quad \text { and } 2 n \equiv 1 \bmod 3 \\
n+5, \quad r=n+1=\Delta, \quad n \geq 5 \quad \text { and } 2 n \equiv 2 \bmod 3
\end{array}\right.
$$

Proof. We divide the proof into some cases.
Case 1 : For $1 \leq r \leq n-1$
The $r$ - dynamic $(n)-$ coloring is as follows:
$\left|N\left(e_{i}\right)\right|=d\left(e_{i}\right)=n-1$,
$\left|N\left(s_{i}\right)\right|=d\left(s_{i}\right)=3=\delta$.
Now an easy check shows that the $r$-adjacency condition is fulfilled.
Hence, $\chi_{r}\left(L\left(G_{n}\right)\right)=n$, for $1 \leq r \leq n-1$

## Case 2: For $r=n$ and $n \not \equiv 1 \bmod 3$

The $r$ - dynamic $(n+2)-$ coloring is as follows:

- Color vertex $e_{i}$ with color $i, 1 \leq i \leq n$.

Let us notice that vertices adjacent to each vertex $e_{i}$ mustbe colored with $r=n$ different colors. After this step each vertex $e_{i}$ has $n-1$ neighbours in different colors and exactly its two neighbours are uncolored: $s_{i-1}, s_{i}$.

We have to color them with atleast one new color to vertex $s_{i}$ to fulfill $r$ adjacenct condition for vertex $s_{i}$. so $\chi_{r}\left(L\left(G_{n}\right)\right) \geq n+2$.

To color vertices $s_{i}, 1 \leq i \leq n$.
Now the number of vertices $s_{i}$, forming a cycle $C_{2 n}$, is not divisibly by 3 , so color the vertices $s_{1}, s_{4}, s_{7}, s_{10}, \cdots s_{2 n-2}$ with color $n+1$.

Now another neighbour of $e_{1}$ has uncolored. So we have to assign one of the allowed colors $c_{1}, c_{2}, c_{3}, \cdots, c_{n}$ to vertex $s_{2 n}$.

Next, the two neighbours of $e_{2}$ are uncolored. We have to color them with atleast one new color to vertex $s_{2}$ to fulfill $r$ - adjacent condition for vertex $e_{i}$.

- color the vertices $s_{2}, s_{5}, s_{8}, \cdots s_{2 n-1}$ with color $n+2$.

Now the neighbours of $e_{i}$ has atleast $n$ colors.
Now $s_{3}, s_{6}, s_{9}, s_{12}, \cdots s_{2 n}$ vertices get any one of the allowed colors $c_{1}, c_{2}, c_{3}, \cdots c_{n}$. Now an easy check shows that the $r$-adjacency condition is fulfilled.
Hence, $\chi_{r}\left(L\left(G_{n}\right)\right)=n+2$, for $r=n$ and $n \not \equiv 1 \bmod 3$

Case 3 : For $r=n$ and $n \equiv 1 \bmod 3$
The $r$ - dynamic $(n+3)-$ coloring is as follows:

- Color vertex $e_{i}$ with color $i, 1 \leq i \leq n$.

Let us notice that vertices adjacent to each vertex $e_{i}$ mustbe colored with $r=n$ different colors. After this step each vertex $e_{i}$ has $n-1$ neighbours in different colors and exactly its two neighbours are uncolored: $s_{i-1}, s_{i}$.

We have to color them with atleast one new color to vertex $s_{i}$ to fulfill $r$ adjacenct condition for vertex $s_{i}$. so $\chi_{r}\left(L\left(G_{n}\right)\right) \geq n+2$.

To color vertices $s_{i}, 1 \leq i \leq n$.
Now the number of vertices $s_{i}$, forming a cycle $C_{2 n}$, is not divisibly by 3 , so color the vertices $s_{1}, s_{4}, s_{7}, s_{10}, \cdots s_{2 n-4}$ with color $n+1$.

Now another neighbour of $e_{1}$ has uncolored. So we have to assign one of the allowed colors $1,2,3, \cdots n$ to vertex $s_{2 n}$.

Next, the two neighbours of $e_{2}$ are uncolored. We have to color them with atleast one new color to vertex $s_{2}$ to fulfill $r$ - adjacent condition for vertex $e_{i}$.

- color the vertices $s_{2}, s_{5}, s_{8}, \cdots s_{2 n-3}$ with color $n+2$.

But the neighbours of $e_{n}$ having only $n-1$ colors. So we have to assign any one of the new color to the vertices $s_{2 n-1}, s_{2 n-2}$.

Suppose to assign color $n+3$ to $s_{2 n-2}$, next assign the uncolored vertices to the any one of the allowed colors $1,2 \cdots n$ to fulfill $r$-adjacent condition for vertex $e_{i}$.

Now an easy check shows that the $r$-adjacency condition is fulfilled.
Hence, $\chi_{r}\left(L\left(G_{n}\right)\right)=n+3$, for $r=n$ and $n \equiv 1 \bmod 3$

Case $4: r=n+1=\Delta$ and $2 n \equiv 0 \bmod 3$
The $r$ dynamic $(n+3)$-coloring is as follows:

- color the vertex $e_{i}$ with color $i, 1 \leq i \leq n$.

It is clear that to color $2 n$ remaining vertices: $s_{i}$ we have to use colors $n, \cdots \chi_{r}$.
we have to still take care of the $r$ - adjacency condition for all vertices. The $r$ - adjacency condition for vertices $s_{i}, 1 \leq i \leq n$, we must use atleast two new colors to vertex $s_{i}$. So $\chi_{r}\left(L\left(G_{n}\right)\right) \geq n+3$.

- Color the vertices $s_{1}, s_{4}, s_{7}, s_{10}, \cdots s_{2 n-2}$ with color $n+1$.
- color the vertices $s_{3}, s_{6}, s_{9}, s_{12}, \cdots s_{2 n}$ with new color $n+2$.

Now the vertex $s_{2}$ is uncolored. So we have to assign the new color $n+3$ to the vertices $s_{2}, s_{5}, s_{8}, \cdots s_{2 n-1}$.

Now an easy check shows that the $r$-adjacency condition is fulfilled.
Hence, $\chi_{r}\left(L\left(G_{n}\right)\right)=n+3$, for $r=n+1=\Delta$ and $2 n \equiv 0 \bmod 3$.

Case 5 : $r=n+1=\Delta$ and $2 n \equiv 1 \bmod 3$
The $r$ dynamic $(n+4)$-coloring is as follows:

- color the vertex $e_{i}$ with color $i, 1 \leq i \leq n$.

It is clear that to color $2 n$ remaining vertices: $s_{i}$ we have to use colors $n, \cdots \chi_{r}$.
we have to still take care of the $r$ - adjacency condition for all vertices. The $r$ - adjacency condition for vertices $s_{i}, 1 \leq i \leq n$, we must use atleast two new colors to vertex $s_{i}$.

- Color the vertices $s_{1}, s_{4}, s_{7}, s_{10}, \cdots s_{2 n-3}$ with color $n+1$.
- color the vertices $s_{2 n}, s_{3}, s_{6}, s_{9}, s_{12}, \cdots s_{2 n-4}$ with new color $n+2$.

Now the vertices $s_{2}, s_{5}, s_{8}, \cdots s_{2 n-2}, s_{2 n-1}$ are uncolored.

- color the vertices $s_{2}, s_{5}, s_{8}, \cdots s_{2 n-1}$ with the color $n+3$.

Now $s_{2 n-2}$ is uncolored. So we have to assign the new color $n+4$ to $s_{2 n-2}$.

Now an easy check shows that the $r$-adjacency condition is fulfilled. Hence, $\chi_{r}\left(L\left(G_{n}\right)\right)=n+4$, for $r=n+1=\Delta$ and $2 n \equiv 1 \bmod 3$.

Case $6: r=n+1=\Delta$ and $2 n \equiv 2 \bmod 3$
The $r$ dynamic $(n+5)$-coloring is as follows:

- color the vertex $e_{i}$ with color $i, 1 \leq i \leq n$.

It is clear that to color $2 n$ remaining vertices: $s_{i}$ we have to use colors $n, \cdots \chi_{r}$.
we have to still take care of the $r$ - adjacency condition for all vertices. The $r$ - adjacency condition for vertices $s_{i}, 1 \leq i \leq n$, we must use atleast two new colors to vertex $s_{i}$.

- Color the vertices $s_{1}, s_{4}, s_{7}, s_{10}, \cdots s_{2 n-4}$ with color $n+1$.
- color the vertices $s_{2 n}, s_{3}, s_{6}, s_{9}, s_{12}, \cdots s_{2 n-5}$ with new color $n+2$.
- color the vertices $s_{2}, s_{5}, s_{8}, \cdots s_{2 n-3}$ with the color $n+3$.

Now $s_{2 n-2}, s_{2 n-1}$ are uncolored.
So we have to assign the new color $n+4$ to $s_{2 n-2}$ and to assign the new color $n+5$ to $s_{2 n-1}$.

Now an easy check shows that the $r$-adjacency condition is fulfilled.
Hence, $\chi_{r}\left(L\left(G_{n}\right)\right)=n+5$, for $r=n+1=\Delta$ and $2 n \equiv 2 \bmod 3$.

In all cases the order of the assigned colors is important. One can verify that the adjacency and $r$-adjacency conditions are fulfilled.

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