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Global neighbourhood domination

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Abstract

A subset D of vertices of a graph G is called a global neighbourhood dominating set(gnd - set) if D is a dominating set for both G and G^N , where G^N is the neighbourhood graph of G. The global neighbourhood domination number(gnd - number) is the minimum cardinality of a global neighbourhood dominating set of G and is denoted by $\gamma_{gn}(G)$. In this paper sharp bounds for γ_{gn} , are supplied for graphs whose girth is greater than three. Exact values of this number for paths and cycles are presented as well. The characterization result for a subset of the vertex set of G to be a global neighbourhood dominating set for G is given and also characterized the graphs of order n having gnd numbers 1, 2, n - 1, n - 2, n.

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1. Introduction & Preliminaries

Domination is an active subject in graph theory, and has numerous applications to distributed computing, the web graph and adhoc networks. For a comprehensive introduction to theoretical and applied facets of domination in graphs the reader is directed to the book [4].

A set D of vertices is called a dominating set of G if each vertex not in D is joined to some vertex in D. The domination number $\gamma(G)$ is the minimum cardinality of the dominating set of G [4].

Many variants of the domination number have been studied. For instance a dominating set S of a graph G is called a restrained dominating set if every vertex in V - D is adjacent to a vertex in D as well as another vertex in V - D. The restrained domination number of G, denoted by $\gamma_r(G)$ is the smallest cardinality of the restrained dominating set of G [3]. A set D is called a global dominating set of G if D is a dominating set of both Gand its complement G^c . The global domination number of G, denoted by $\gamma_g(G)$ is the smallest cardinality of the global dominating set of G [6]. A dominating set D of connected graph G is called a connected dominating set G if the induced subgraph < D > is connected. The connected domination number of G, denoted by $\gamma_c(G)$ is the smallest cardinality of the connected dominating set of G [7]. A dominating set D of connected graph G is called an independent dominating set of G if the induced subgraph < D > is a null graph [4].

G be a connected graph, then the Neighbourhood Graph of G is denoted by N(G) (or) G^N and it has the same vertex set as that of G and edge set being $\{uv/u, v \in V(G), there is w \in V(G) \text{ such that } uw, wv \in E(G)\}$ [2].

Recently we have introduced a new type of graph known as semi complete graph. Let G be a connected graph, then G is said to be semi complete if any pair of vertices in G have a common neighbour. The necessary and sufficient condition for a connected graph to be semi complete is any pair of vertices lie on the same triangle or lie on two different triangles having a common vertex [5].

In the present paper, we introduce a new graph parameter, the global neighbourhood domination number, for a connected graph G. We call $D \subseteq V(G)$ a global neighbourhood dominating set (gnd - set) of G if D is a dominating set for both G, G^N . The global neighbourhood domination number is the minimum cardinality of a global neighbourhood dominating set of G and is denoted by $\gamma_{qn}(G)$.

Example. Suppose G is a graph representing a network of roads linking various locations. Some essential goods are being supplied to these locations from supplying stations. It may happen that these links(edges of G) may be broken for some reason or the other. So we have to think of maintaining the supply of goods to various locations uninterrupted through secret links(edges in the neighbourhood graph of G). As the neighbourhood graph of G is a spanning subgraph of G^c , the construction (maintainance) cost of secret links can be minimized, when compared with the complementary graph of G. The global neighbourhood domination number will be the minimum number of supplying stations needed to accomplish the task of supplying the goods uninterruptedly.

All graphs considered in this paper are simple, finite, undirected and connected. For all graph theoretic terminology not defined here, the reader is referred to [1].

In section [2], sharp bounds for γ_{gn} are supplied for the graphs whose girth is greater than three. In section [3], we have given a characterization result for a proper subset of the vertex set of G to be a gnd - set of G and also characterized the graphs whose gnd - numbers are 1, 2, n, n - 1, n - 2.

2. Bounds for the global neighbourhood domination number

In this section, we obtain some bounds for the gnd - numbers of graphs whose girth is greater than three.

Theorem 2.1. If G is a triangle free graph, then

$$\frac{2e - n(n-3)}{2} \le \gamma_{gn}(G) \le n - \Delta(G) + 1.$$

Proof: Let D be a minimum gnd - set of G. By hypothesis every vertex in V - D is non adjacent with atleast one vertex in D. Otherwise we get a contradiction to that D is a gnd - set for G.

$$\Rightarrow e \leq \frac{n(n-1)}{2} - [n - \gamma_{gn}(G)]$$
$$\Rightarrow \frac{2e - n(n-3)}{2} \leq \gamma_{gn}(G) \qquad \to (1)$$

Suppose that $d_G(v) = \Delta(G)$ for some v in V(G).

Let $v_1, v_2, ..., v_{\Delta(G)}$ be the neighbours of v in G. Then

 $[V - \{v_1, v_2, ..., v_{\Delta(G)}\}] \bigcup \{v_i : i \text{ is one of } 1, 2, ..., \Delta(G)\}$ is a gnd - set of G and its cardinality is $n - \Delta(G) + 1$.

$$\Rightarrow \gamma_{gn}(G) \le n - \Delta(G) + 1 \qquad \to (2)$$

From (1) and (2)

$$\frac{2e - n(n-3)}{2} \le \gamma_{gn}(G) \le n - \Delta(G) + 1.$$

Furthermore the lower bound is attained in the case of C_4 and upper bound is attained in the case of P_3 . Hence the bounds are sharp.

Note: The upper bound holds good for any connected graph G.

1. $\gamma_{gn}(K_n) = 1; n \ge 3$ 2. $\gamma_{gn}(S_n) = 2; n \ge 3$ 3. $\gamma_{gn}(K_{m,n}) = 2; m + n \ge 3$ 4. $\gamma_{gn}(P_n) = [\frac{n}{3}]$; n = 3m + 1

$$= \left[\frac{n}{3}\right] + 2; \ n = 3m, 3m + 2$$

Here $n \geq 4$.

5.
$$\gamma_{gn}(C_n) = [\frac{n}{3}]$$
; $n = 3m$

$$= \left[\frac{n}{3}\right] + 1; \quad n = 3m + 1, 3m + 2$$

6.
$$\gamma_{qn}(C_n O K_2) = n$$

 $G = P_n (n \ge 4)$. Then there is an independent gnd - set for G iff n = 3m + 1.

 $G = C_n (n \ge 3)$. Then there is an independent gnd - set for G iff n = 3m.

 $G = P_n(n \ge 3)$. Then $\gamma_{gn}(G) = n - 2$ iff n = 4, 5.

 $G = C_n (n \ge 3)$. Then $\gamma_{gn}(G) = n - 2$ iff n = 3, 4, 5.

If T is a tree of order $n \ge 3$, then $\gamma_{gn}(G) = 2$ iff T is obtained from P_3 or P_4 by adding zero or more leaves to the stems of the path.

Theorem 2.2. G be a connected graph and D be a minimum dominating set of G. If there is a vertex v in V - D such that v is adjacent to all the vertices in D, then $\gamma_{gn}(G) \leq 1 + \gamma(G)$.

Proof: Assume that $D \subseteq N_G(v)$ for some $v \in V$. The proof follows from the fact that $D \bigcup \{v\}$ is a gnd - set of G.

Theorem 2.3. D be a minimum dominating set of G. Then $\gamma_{gn}(G) = 1 + \gamma(G)$ iff there is a vertex v in V - D satisfying:

 $(i)N(v) \subset D$, each of the vertices in N(v) is isolated in $\langle D \rangle$.

(ii) $v_1 \in V - D(v \neq v_1)$ satisfies (i) then $N(v) \cap N(v_1) \neq \phi$.

Proof: Assume that $\gamma_{gn}(G) = 1 + \gamma(G)$. Then there is a vertex v in V - D satisfying (i) and (ii), otherwise $\gamma_{gn}(G) = \gamma(G)$ which is a contradiction.

Assume that the converse holds. Then $D \cup \{v\}$ is a gnd - set in G and D is not a gnd - set in G. Thus $D \cup \{v\}$ is a minimum gnd - set in G.

Hence $\gamma_{gn}(G) = |D \bigcup \{v\}| = \gamma(G) + 1.$

Theorem 2.4. G be a connected graph, then $\gamma(G) \leq \gamma_{gn}(G) \leq \gamma_c(G)$.

Proof: Clearly $\gamma(G) \leq \gamma_{gn}(G)$. Since any connected dominating set for G is a gnd - set for $G, \gamma_{gn}(G) \leq \gamma_c(G)$. Hence $\gamma(G) \leq \gamma_{gn}(G) \leq \gamma_c(G)$.

Theorem 2.5. G be a connected graph with g(G) > 3, then $\gamma_g(G) \leq \gamma_{qn}(G)$.

Proof: By hypothesis, every gnd - set is a global dominating set in G. Hence $\gamma_g(G) \leq \gamma_{gn}(G)$.

Note: Under the hypothesis given in the Theorem(2.10) and Theorem(2.9), we have

$$\gamma(G) \le \gamma_g(G) \le \gamma_{gn}(G) \le \gamma_c(G).$$

Theorem 2.6. G be a connected graph. Then $G = G^{NN}$ iff

(i) Each edge in G lies on C_3 or C_5 .

(ii) There is no path of length four between any pair of non adjacent vertices in G.

Proof: Assume that $G = G^{NN}$.

(i) Let v_1, v_2 be an arbitrary edge in G, then by our assumption v_1v_2 is an edge in G^{NN} .

Suppose $v_1v_2 \in G^N$. Since $v_1v_2 \in E(G), v_1v_2$ lies on a cycle C_3 in G.

Suppose $v_1v_2 \notin G^N$. Since $v_1v_2 \in G^{NN}$, there is a v_3 in V(G) such that $\langle v_1v_3v_2 \rangle$ is a path in G^N . This implies $v_1v_3, v_3v_2 \in E(G^N)$. So there is a path of length four from v_1 to v_2 in G. Thus $P \bigcup \{v_1v_2\}$ is a 5 - cycle in G. Therefore v_1v_2 lies on C_5 .

Hence each edge in G lies on C_3 or C_5 .

(ii) If there is a path of length four between any pair of non adjacent vertices in G, then there is an edge in G^{NN} which is not in G. Hence $G \neq G^{NN}$, which is a contradiction.

Assume that the converse holds.

Let v_1v_2 be an arbitrary edge in G. Then by (i) of our assumption v_1v_2 lies on C_3 or C_5 . In either case v_1v_2 is an edge in G^{NN} . Hence $G \subseteq G^{NN}$.

Let $v_1v_2 \in E(G^{NN})$. Suppose that $v_1v_2 \notin E(G)$. In either case $v_1v_2 \in E(G^N)$ or $v_1v_2 \notin E(G^N)$ there is a path of length four from v_1 to v_2 in G, which is a contradiction. Thus $v_1v_2 \in E(G)$. Thus $G^{NN} \subseteq G$.

Hence $G^{NN} = G$.

If G is a graph satisfying conditions (i) and (ii) of Theorem (2.12), then

$$\frac{\gamma + \overline{\gamma}}{2} \le \gamma_{gn} \le \gamma + \overline{\gamma}.$$

Proof: Under the given hypothesis by Theorem(2.12), $\gamma_{gn} = \overline{\gamma}_{gn}$.

We have $\gamma \leq \gamma_{gn}, \overline{\gamma} \leq \overline{\gamma}_{qn} = \gamma_{gn}$

$$\Rightarrow \quad \gamma + \overline{\gamma} \le 2\gamma_{gn}$$
$$\Rightarrow \quad \frac{\gamma + \overline{\gamma}}{2} \le \gamma_{gn} \to (1)$$

Clearly,

$$\gamma_{gn} \le \gamma + \overline{\gamma} \qquad \to (2)$$

From (1) and (2),

$$\frac{\gamma + \overline{\gamma}}{2} \le \gamma_{gn} \le \gamma + \overline{\gamma}.$$

Here $\gamma = \gamma(G), \gamma_{gn} = \gamma_{gn}(G), \overline{\gamma} = \gamma_{gn}(G^N).$

3. Characterization and Other Relevant Results.

In this section we have given the characterization for a proper subset of the vertex set of a graph to be a gnd - set.

Theorem 3.1. (Characterization Result) *G* be a connected graph. $D \subset V$ is a gnd - set of *G* iff each vertex in V - D lies on an edge whose end points are totally dominated by the vertices in *D*.

Proof: Assume that D is a gnd - set for G. Let $v_1 \in V - D$.

Since D is a dominating set for G there is a $v_2 \in D$ such that $v_1v_2 \in E(G)$.

Since D is a dominating set for G^N there is a $v_3 \in D$ such that $\langle v_1v_4v_3 \rangle$ is a path in G for some $v_4 \in V(G)$.

If $v_4 \neq v_2$, then v_1 lies on the edge v_1v_4 , where v_1 is dominated by v_2 and v_4 is dominated by $v_3(v_2, v_3 \in D)$.

If $v_4 = v_2$, then v_1 lies on the edge v_1v_2 , where v_1 is dominated by v_2 and v_2 is dominated by $v_3(v_2, v_3 \in D)$.

So in either case v_1 lies on the edge whose end points are totally dominated by the vertices in D.

Assume that the converse holds.

Let $v_1 \in V - D$. Then by our assumption there is a $v_2 \in V(G), v_3, v_4 \in D$ such that $v_1v_3, v_2v_4 \in E(G)$.

case:(i) Suppose $v_2 = v_3$.

Then $\langle v_1 v_2 v_4 \rangle$ is a path in G

 $\Rightarrow v_1 v_4 \in E(G^N), v_4 \in D.$

case:(ii) Suppose $v_2 \neq v_3$.

Then $\langle v_3 v_1 v_2 v_4 \rangle$ is a path in G

 $\Rightarrow v_1v_4 \in E(G^N), v_4 \in D.$

Therefore v_1 is dominated by v_3 in G and by v_4 in G^N .

Since v_1 is arbitrary, D is a gnd - set of G.

Theorem 3.2. G be a connected graph. Then G^N is complete iff G is semi-complete.

Proof: Assume that G^N is complete.

Let $v_1, v_2 \in V(G)$. Since G^N is complete, $v_1v_2 \in E(G^N)$. Then there is a v_3 in $V(G) - \{v_1, v_2\}$ such that $\langle v_1v_3v_2 \rangle$ is a path in G.

Hence G is semi complete. Assume that the converse holds.

Let $v, v_2 \in V(G)$. By our assumption, there is a v_3 in G such that $\langle v_1 v_3 v_2 \rangle$ is a path in G

 $\Rightarrow v_1 v_2 \in E(G^N).$

Hence G^N is complete. G be a connected graph. Then $\gamma_{gn}(G) = 1$ iff G is semi-complete and $\gamma(G) = 1$.

Proof: Assume that $\gamma_{gn}(G) = 1$.

$$\Rightarrow \gamma_{(G)} = 1, \gamma_{gn}(G) = 1.$$

Then there is a $v_0 \in V$ such that $D = \{v_0\}$ is a dominating set for G, G^N .

Now we show that G is semi complete. Let $v_1, v_2 \in V(G)$.

case:(i) Suppose $v_1 \neq v_0 \neq v_2$.

subcase:(a) Suppose $v_1v_2 \in E(G)$.

Clearly $\langle v_1 v_0 v_2 \rangle$ is a path in G. Thus $\langle v_1 v_0 v_2 \rangle \bigcup \{v_1 v_2\}$ is a triangle in G.

Hence v_1v_2 lie on the same triangle in G.

subcase:(b) Suppose $v_1v_2 \notin E(G)$.

Clearly $\langle v_1v_0v_2 \rangle$ is a path in G. Since $v_1v_0 \in E(G)$ and $\{v_0\}$ is a dominating set of G^N there is a $v_3 \in V - \{v_0, v_1\}$ such that $\langle v_1v_3v_0v_1 \rangle$ is a triangle in G.

Similarly for v_0v_2 in E(G), there is v_4 in $V - \{v_0, v_2\}$ such that $\langle v_2v_0v_4v_2 \rangle$ is a triangle in G.

Hence v_1, v_2 lie on triangles $\langle v_1v_3v_0v_1 \rangle$, $\langle v_2v_0v_4v_2 \rangle$ respectively and v_0 is the common vertex.

case:(ii) One of v_1, v_2 is v_0 .

Without loss of generality assume that $v_1 = v_0$. Since $\{v_0\}$ is a dominating set for $G, v_0v_2 \in E(G)$. Then $\langle v_0v_2v_3v_0 \rangle$ is a triangle in G for some v_3 in $V - \{v_0, v_2\}$. Thus v_0, v_2 lie on the same triangle in G.

Therefore from case(1) and case(2), any two vertices in G lie on the same triangle or they lie on different triangles having a common vertex.

Hence by the characterization theorem for semi complete graphs, G is semi complete.

Assume that the converse holds.

Since G is semi complete, by Theorem (3.2) G^N is complete.

 $\Rightarrow \gamma(G^N) = 1.$

Hence $\gamma_{gn}(G) = 1$. G be a semi complete graph. Then $\gamma_{gn}(G) = 1$ iff G is a union of triangles having a common vertex.

Proof: Assume that $\gamma_{qn}(G) = 1$

 $\Rightarrow \gamma(G) = 1, \gamma_{gn}(G) = 1.$ Let $D = \{v_0\}$ is a dominating set for G, G^N for some $\{v_0\}$ in V.

Let $\{v_1, v_2 \in E(G)\}.$

Case:1. Suppose $v_1 \neq v_0 \neq v_2$.

Then $\langle v_1 v_0 v_2 v_1 \rangle$ is a triangle in G.

Case:2. One of v_1, v_2 is v_0 .

Without loss of generality assume that $v_1 = v_0$. since v_0v_2 is an edge in G and G is semi complete there is a $v_3 \in G$ such that $\langle v_0v_2v_3v_0 \rangle$ is a triangle in G.

Since v_1v_2 is an arbitrary edge, each edge lies on a triangle having v_0 as the common vertex.

Assume that the converse holds.

By our assumption $\gamma(G) = 1$. Then by Theorem(3.3) $\gamma_{qn}(G) = 1$.

Theorem 3.3. *G* be a connected graph and *H* be a spanning subgraph of *G*, then $\gamma_{gn}(G) \leq \gamma_{gn}(H)$.

Theorem 3.4. G be a connected graph, then $\gamma_{qn}(G) = n$ iff $G = K_2$.

Proof: Assume that $\gamma_{gn}(G) = n$.

Suppose that $diam(G) \geq 2$. Let $\langle v_1v_2...v_{k-1}v_k \rangle$ $(k \geq 3)$ be the diammetral path in G. Then $V - \{v_1\}$ or $V - \{v_k\}$ is a gnd - set in G, which is a contradiction to our assumption. So, diam(G) = 1. This implies $G = K_n (n \geq 2)$.

By Corollary.3.3. n = 2 (i.e) $G = K_2$. The converse part is clear.

Theorem 3.5. G be a connected graph, then $\gamma_{gn}(G) = n - 1$ iff $G = P_3$.

Proof: Assume that $\gamma_{gn}(G) = n - 1$.

Suppose that $diam(G) \geq 3$. Let $\langle v_1v_2...v_{k-1}v_k \rangle$ $(k \geq 4)$ be the diammetral path in G. Then $V - \{v_1, v_k\}$ is a gnd - set in G, which is a contradiction to our assumption. Hence $diam(G) \leq 2$. By the above Theorem and by our assumption $diam(G) \neq 1$. So diam(G) = 2.

If $G = P_3$, then $\gamma_{qn}(G) = n - 1$.

Suppose $G \neq P_3$. Form a spanning tree G' of G. Clearly $diam(G') \geq 3$. This implies $\gamma_{gn}(G) < n-1$, a contradiction. Hence our supposition is false. The converse part is clear.

Theorem 3.6. G be a connected graph of order $n \ge 4$, then $\gamma_{gn}(G) = n-2$ iff $G = C_n (3 \le n \le 5)$ or $G = P_n (n = 4, 5)$ or $G = S_3$ or G is isomorphic to

Proof: Assume that $\gamma_{gn}(G) = n - 2$.

Suppose $diam(G) \ge 5$. Form a spanning tree G' of G. Clearly $diam(G') \ge 5$.

If G' has more than two pendant vertices, then by Theorem.3.5. $V(G') - \{pendant vertices in G'\}$ is a gnd - set of G of cardinality less than n-2. Hence a contradiction to our assumption. If G' has exactly two pendant vertices, then $G' = \langle v_1 v_2 ... v_{k-1}, v_k \rangle$ $(k \ge 6)$. By Theorem 3.5. $V - \{v_1, v_4, v_{k-3}, v_k\}$ is a gnd - set of G of cardinality less than n - 2, a contradiction.

Hence $diam(G) \le 4$. Suppose diam(G) = 4. Let $P_5 = \langle v_1 v_2 v_3 v_4 v_5 \rangle$ be a diammetral path in G.

Case:1. $V(G) = V(P_5)$.

We have two possibilities $G = P_5$ or $G \neq P_5$. If $G = P_5$, then $\gamma_{gn}(G) = n-2$.

Suppose $G \neq P_5$.

If g(G) = 3, then $\gamma_{qn}(G) = 2 < n-2$ which is a contradiction.

If g(G) = 4, then G is isomorphic to

In either case $\gamma_{gn}(G) = 2 < n - 2$.

If g(G) = 5, then $G \cong C_5$. This implies $\gamma_{gn}(G) = n - 2$.

Case:2. $V(G) \neq V(P_5)$.

Form a spanning tree G' of G. Clearly $diam(G') \ge 4$ and G' has at least three pendant vertices, which implies $\gamma_{gn}(G) < n-2$, a contradiction.

Suppose diam(G) = 3.

Let $P_4 = \langle v_1 v_2 v_3 v_4 \rangle$ be a diammetral path in G.

Case:1. $V(G) = V(P_4)$.

We have two possibilities $G = P_4$ or $G \neq P_4$. If $G = P_4$, then $\gamma_{gn}(G) = n-2$.

Suppose $G \neq P_4$. Clearly $g(G) \leq 4$.

If g(G) = 3, then G is a union of triangles having a common vertex or G is isomorphic to H. In the former case, by Corollary.3.4. we get $\gamma_{gn}(G) = 1 < n-2$ a contradiction. In the later case we have $\gamma_{gn}(G) = n-2$.

If g(G) = 4, then $\gamma_{qn}(G) = n - 2$ and $G = C_4$.

Case:2. $V(G) \neq V(P_4)$. Then as in the case:2 of diam(G) = 4 we get a contradiction.

Suppose diam(G) = 2.

Let $P_3 = \langle v_1 v_2 v_3 \rangle$ be a diammetral path in G.

Case:1. $V(G) = V(P_3)$.

We have two possibilities $G = P_3$ or $G \neq P_3$. If $G = P_3$, then $\gamma_{gn}(G) > n-2$ a contradiction. If $G \neq P_3$, then $G = C_3$. Thus $\gamma_{gn}(G) = n-2$.

Case:2. $V(G) \neq V(P_3)$.

Form a spanning tree G' of G. Clearly $diam(G') \ge 2$.

If G' has more than one internal vertex then, $V(G) - \{\text{pendant vertices of } G'\}$ is a gnd - set of G'. Hence a gnd - set of G of cardinality less than n - 2 a contradiction.

If G' has exactly one internal vertex $((i.e)v_2)$, then $G' \cong S_n (n \ge 2)$. If $n \ge 4$, then $\gamma_{gn}(G) < n-2$ a contradiction. If $G' = S_2$ again we get a contradiction. So $G' = S_3$. This implies $G = S_3$ or G is a union of triangles having a common vertex or G is isomorphic to H. If G is a union of triangles having a common vertex then by Corollary.3.4. we get a contradiction to $\gamma_{gn}(G) = n-2$. In the remaining cases the equality holds good.

Suppose diam(G) = 2. Hence $G = K_n (n \ge 2)$. If n = 3 we have $\gamma_{an}(G) = n - 2$, otherwise we get a contradiction.

The converse part is clear.

Theorem 3.7. *G* be a connected graph with $g(G) \neq 3$. Then $\gamma_{gn}(G) = 2$ iff there is an edge uv in *G* with $N(u) \bigcup N(v) = V$.

Proof: Assume that $\gamma_{gn}(G) = 2$. Then there is a $D = \{u, v\}$, a gnd - set of G. Clearly $uv \in E(G)$ (if not there is a $w \in V$ which is not dominated by either u or v in G^N , a contradiction). Also since D is a dominating set for $G N(u) \bigcup N(v) = V$.

Assume that the converse holds. Let $D = \{u, v\}$. By our assumption D is a dominating set of G. Since $g(G) \neq 3, D$ is a dominating set of G^N . So D is a gnd - set for G. Since $g(G) \neq 3$, G cannot be semi complete. Then by the Corollary.3.3., D is a gnd - set of minimum cardinality. Hence $\gamma_{gn}(G) = 2$.

Theorem 3.8. *G* be a semi complete graph and $D \subset V$. Then *D* is a dominating set for *G* iff *D* is a gnd - set for *G*.

Proof: Suppose that D is a dominating set for G. Let $v_1 \in V - D$. Then there is a $v_2 \in D$ such that v_1v_2 is an edge in G. Since G is semi complete there is $v_3 \in D - \{v_1, v_2\}$ such that $v_1v_2v_3v_2$ is a triangle in G. Thus v_1v_3, v_3v_2 are edges in G

 $\Rightarrow v_1v_2 \in E(G), v_2 \in D$ Hence D is a gnd - set for G. The converse part is clear.

Theorem 3.9. *G* be a connected graph and $D \subset V$ is independent. Then *D* is a gnd - set for *G* iff *D* is a restrained dominating set for *G*.

Proof: Assume that D is a restrained dominating set for G.

Let $v_1 \in V - D$. By our assumption there is v_2 in $V - D, v_3$ in D such that v_1v_2, v_1v_3 are in E(G). Since $v_2 \in V - D$ and D is a dominating set for G there is a v_4 in D such that $v_2v_4 \in E(G)$.

If v_1v_4 is in G, then $\langle v_1v_2v_4v_1 \rangle$ is a path in $G \Rightarrow v_1v_4 \in E(G^N)$.

If v_1v_4 is not in G, since v_2 is a common neighbour to $v_1, v_4 \Rightarrow v_1v_4 \in E(G^N)$.

D is a gnd - set of G.

Conversely assume that D is a gnd - set for G.

Let $v_1 \in V - D$, by our assumption there is a v_2 in D such that v_1v_2 is in G. Since D is a dominating set for G^N there is a v_3 in V - D such that v_1v_3 is an edge in G. Otherwise we get a contradiction to that D is independent.

Since $v_1 \in V - D$ is arbitrary, D is a restrained dominating set for G. If $G \neq K_{m,n}$ and G is bipartite, then none of the partite sets of G can form a gnd - set for G.

Proof: Let V_1, V_2 be the partite sets for G.

Assume that V_1 is a gnd - set for G.

Let $v_1 \in V - V_1 \Rightarrow v_1 \in V_2$. Since V_1 is a gnd - set for G, by the characterization theorem for gnd - set there is a v_2 in G and v_3, v_4 in V_1 such that $\langle v_3 v_1 v_2 v_4 \rangle$ is a path in G.

case:(i) Suppose $v_2 = v_3$.

Then $\langle v_1v_3v_4 \rangle$ is a path in G, where $v_3, v_4 \in V_1$ which is a contradiction to that G is bipartite.

case:(ii) Suppose $v_2 \neq v_3$.

Then $v_2 \in V_1$ (or) V_2 we get a contradiction to that G is bipartite. Hence V_1 is not a gnd - set for G.

Similarly V_2 is not a dominating set for G.

 $G = C_n OK_2$. If $D \subset V$ is a minimum gnd - set of G, then D is independent iff $D = \{v \in V(G) : d_G(v) = 1\}$.

Proof: Assume that D is independent.

Suppose that $v_1 \in D$ and $d_G(v) \neq 1$. Since D is a minimum independent gnd - set for G there is a $v_2 \in V - D$ such that v_1v_2 is an edge in G and $d_G(v_2) = 1$.

Also there is an edge in G such that v_1v_3 is an edge in G. Hence a contradiction. Thus $v_1 \in D$, then $d_G(v_1) = 1$.

Hence $D = \{v \in V(G) : d_G(v) = 1\}.$

The converse part is clear.

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